

Scala Collections Performance

GS.com/Engineering
March, 2015
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Who am I?

- Engineer at Goldman Sachs
- Technical lead for **GS Collections**
 - A replacement for the **Java Collections Framework**

Goals

- **Scala** programs ought to perform as well as **Java** programs, but **Scala** is a little slower

From `RedBlackTree.scala`

```
/*  
 * Forcing direct fields access using the @inline annotation  
 * helps speed up various operations (especially  
 * smallest/greatest and update/delete).  
 *  
 * Unfortunately the direct field access is not guaranteed to  
 * work (but works on the current implementation of the Scala  
 * compiler).  
 *  
 * An alternative is to implement the these classes using plain  
 * old Java code...  
 */
```

Goals

- **Scala** programs ought to perform as well as **Java** programs, but **Scala** is a little slower
- Highlight a few performance problems that matter to me
- Suggest fixes, borrowing ideas and technology from **GS Collections**

Goals

GS Collections and Scala's Collections are similar

- Mutable and immutable interfaces with common parent
- Similar iteration patterns at hierarchy root
`Traversable` and `RichIterable`
- Lazy evaluation (`view`, `asLazy`)
- Parallel-lazy evaluation (`par`, `asParallel`)

Agenda

GS Collections and Scala's Collections are different

- Persistent data structures
- Hash tables
- Primitive specialization
- Fork-join

Persistent Data Structures

Persistent Data Structures

- Mutable collections are similar in GS Collections and Scala – though we'll look at some differences
- Immutable collections are quite different
- **Scala's** immutable collections are **persistent**
- **GS Collections'** immutable collections are not

Persistent Data Structures

Wikipedia:

In computing, a **persistent data structure** is a data structure that always preserves the previous version of itself when it is modified.

Persistent Data Structures

Wikipedia:

In computing, a **persistent data structure** is a data structure that always preserves the previous version of itself when it is **modified**.

Operations yield new updated structures when they are “modified”

Persistent Data Structures

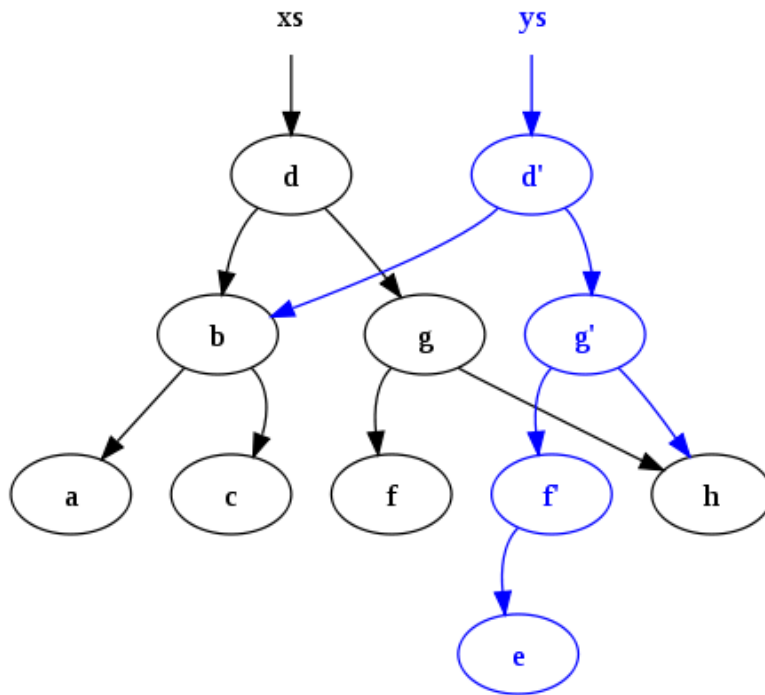
- Typical examples: List and Sorted Set

Persistent Data Structures

- Typical examples: List and **Sorted Set**
- “Add” by constructing $O(\log n)$ new nodes from the leaf to a new root.
- **Scala sorted set** → binary tree
- **GSC immutable sorted set** → covered later

Persistent Data Structures

Wikipedia:



Persistent Data Structures

- Typical examples: **List** and Sorted Set
- “Prepend” by constructing a new head in $O(1)$ time. LIFO structure.
- **Scala immutable list** → linked list or stack
- **GSC immutable list** → array backed

Persistent Data Structures

- Extremely important in purely functional languages
- All collections are immutable
- Must have good runtime complexity
- No one seems to miss them in GS Collections

Persistent Data Structures

- Proposal: Mutable, **Persistent Immutable**, and **Plain Immutable**
- Mutable: Same as always
- Persistent: Use when you want structural sharing
- (Plain) Immutable: Use when you're done mutating or when the data never mutates

Persistent Data Structures

Plain old immutable data structures

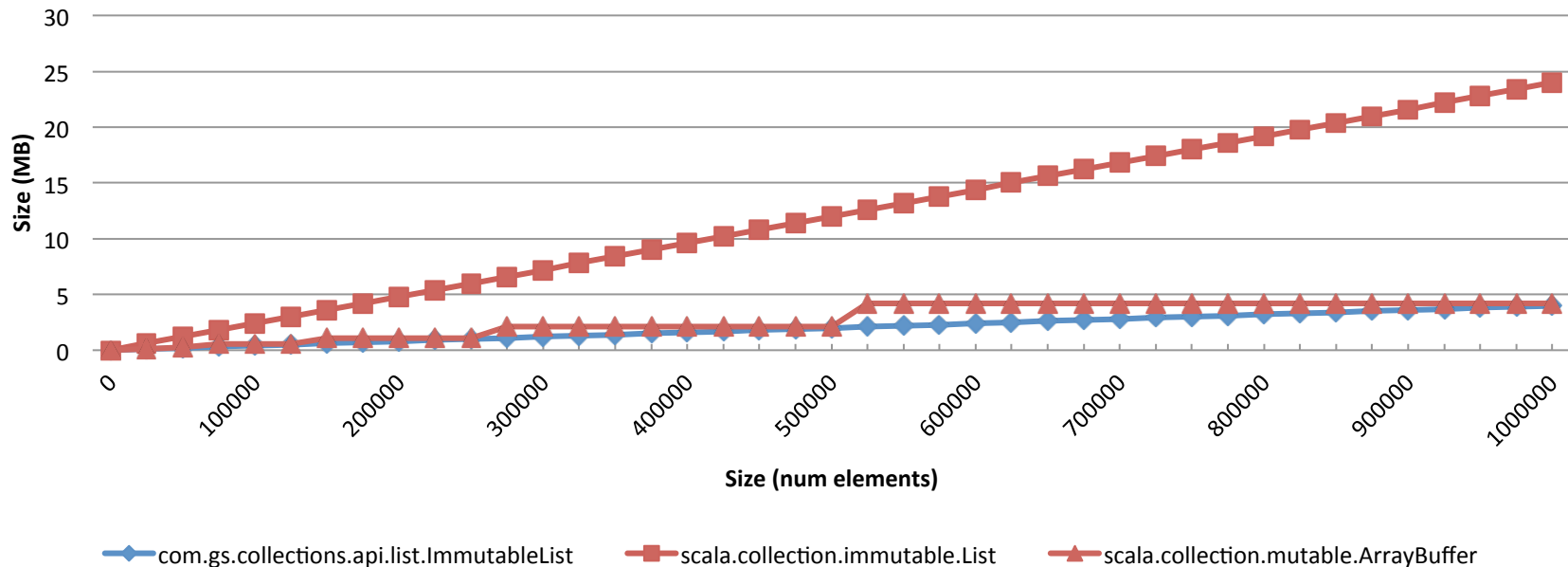
- Not persistent (no structural sharing)
- “Adding” is much slower: $O(n)$
- Speed benefits for everything else
- Huge memory benefits

Persistent Data Structures

- Immutable array-backed list
- Immutable → trimmed array
- No nodes → 1/6 memory
- Array backed → cache locality, should parallelize well

Persistent Data Structures

Lists: Size in MB by number of elements



Measured with Java 8 64-bit and compressed oops

Persistent Data Structures

Performance assumptions:

- Iterating through an array should be faster than iterating through a linked list
- Linked lists won't parallelize well with `.par`
- No surprises in the results – so we'll skip

github.com/goldmansachs/gs-collections/tree/master/jmh-tests

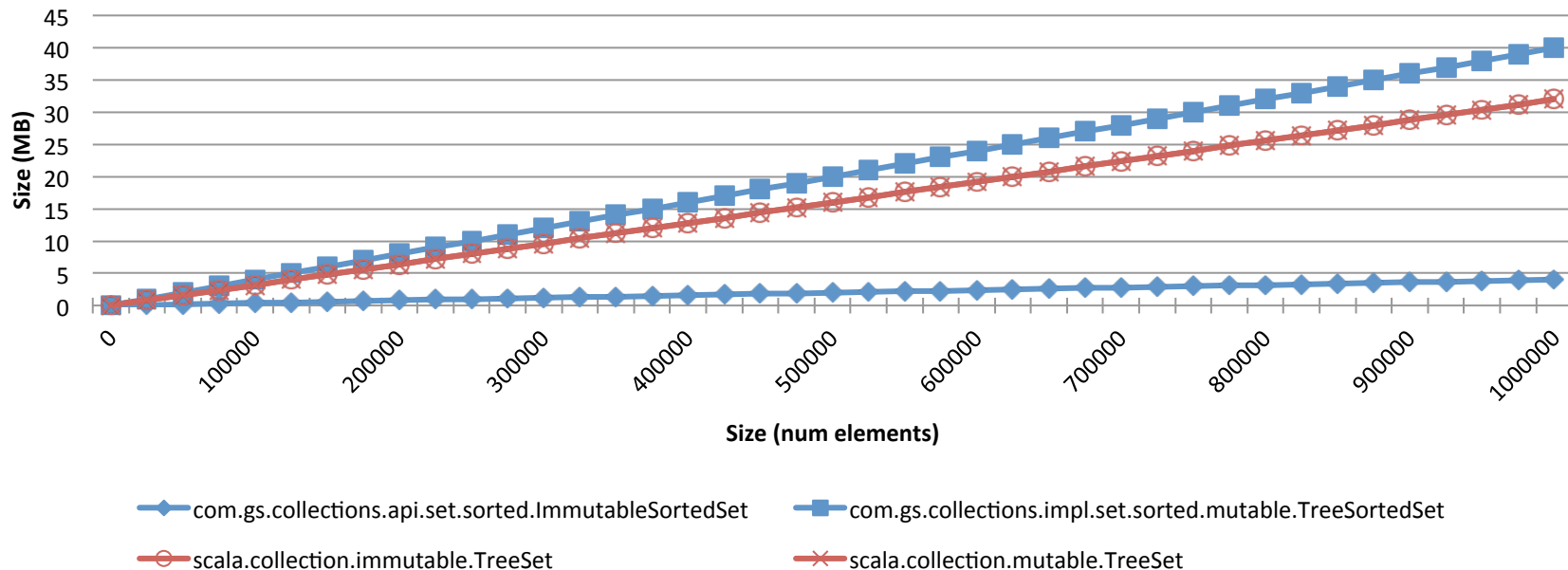
github.com/goldmansachs/gs-collections/tree/master/memory-tests

Persistent Data Structures

- Immutable array-backed sorted set
- Immutable → trimmed, sorted array
- No nodes → ~1/8 memory
- Array backed → cache locality

Persistent Data Structures

Sorted Sets: Size in MB by number of elements

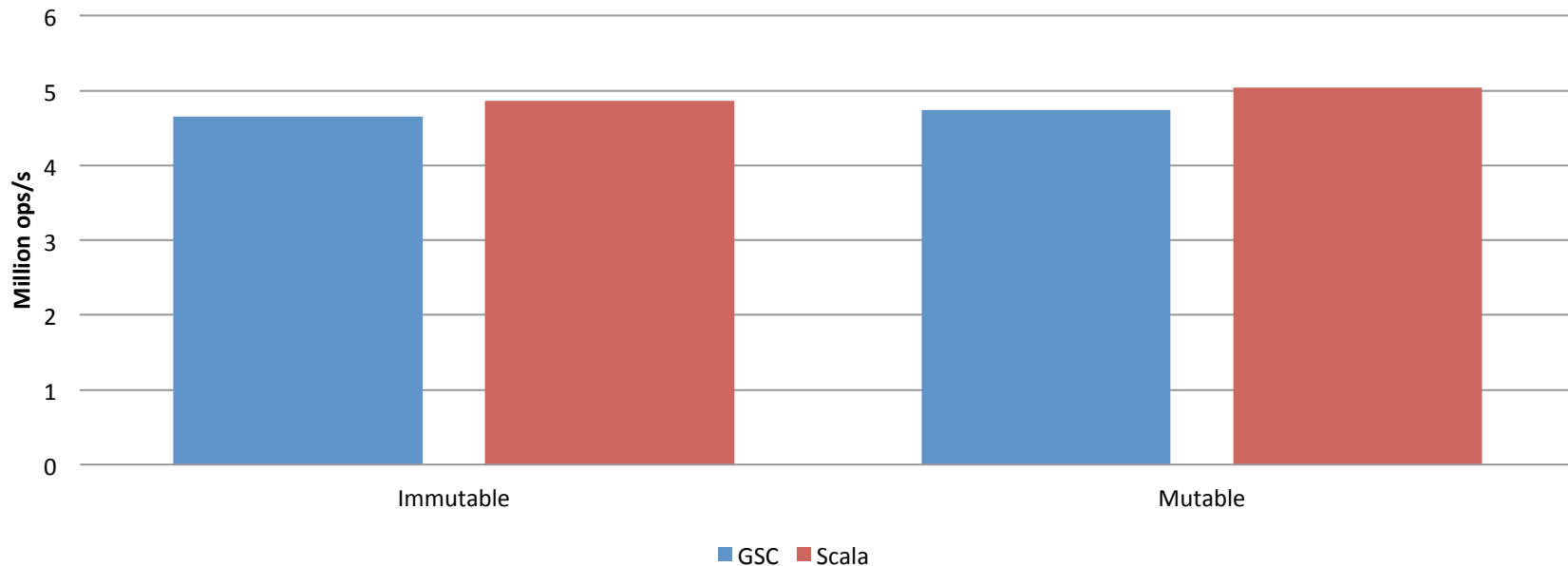


Persistent Data Structures

- Assumptions about contains:
- May be faster when it's a binary search in an array (good cache locality)
- Will be about the same in mutable / immutable tree (all nodes have same structure)

Persistent Data Structures

Sorted Set contains() throughput (higher is better)



Iteration

Persistent Data Structures

```
val scalaImmutable: immutable.TreeSet[Int] =
  immutable.TreeSet.empty[Int] ++ (0 to SIZE)

def serial_immutable_scala(): Unit =
{
  val count: Int = this.scalaImmutable
    .view
    .filter(each => each % 10000 != 0)
    .map(String.valueOf)
    .map(Integer.valueOf)
    .count(each => (each + 1) % 10000 != 0)
  if (count != 999800)
  {
    throw new AssertionError
  }
}
```



Persistent Data Structures

```
val scalaImmutable: immutable.TreeSet[Int] =  
  immutable.TreeSet.empty[Int] ++ (0 to SIZE)  
  
def serial_immutable_scala(): Unit =  
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  if (count != 999800)  
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  }  
}
```

Lazy evaluation so
we don't create
intermediate
collections



Persistent Data Structures

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  if (count != 999800)  
  {  
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  }  
}
```

Lazy evaluation so
we don't create
intermediate
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Terminating in .toList
or .toBuffer would be
dominated by
collection creation



Persistent Data Structures



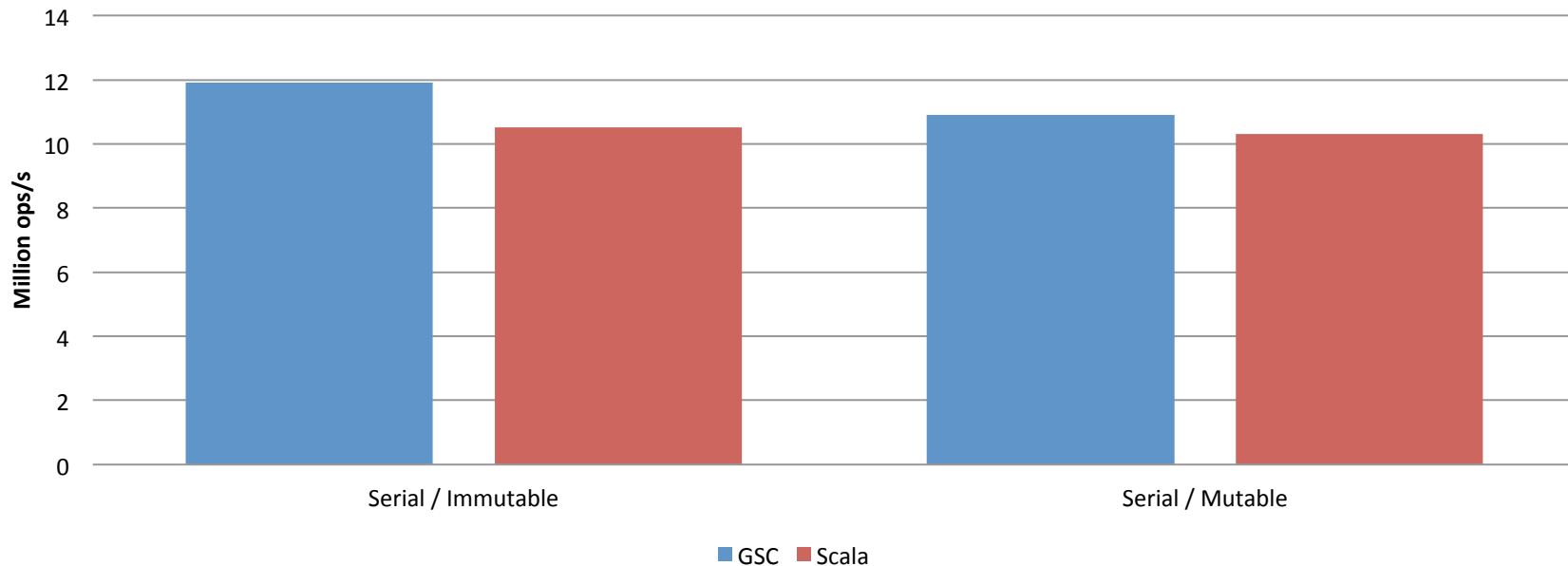
```
private final ImmutableSortedSet<Integer> gscImmutable =
    SortedSets.immutable.withAll(Interval.zeroTo(SIZE));
@Benchmark
public void serial_immutable_gsc()
{
    int count = this.gscImmutable
        .asLazy()
        .select(each -> each % 10_000 != 0)
        .collect(String::valueOf)
        .collect(Integer::valueOf)
        .count(each -> (each + 1) % 10_000 != 0);
    if (count != 999_800)
    {
        throw new AssertionError();
    }
}
```

Persistent Data Structures

- Assumption: Iterating over an array is somewhat faster

Persistent Data Structures

Sorted Set Iteration throughput (higher is better)



Persistent Data Structures

Next up, parallel-lazy evaluation

```
this.scalaImmutable.view →
```

```
this.scalaImmutable.par
```

```
this.gscImmutable.asLazy() →
```

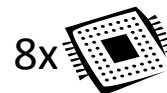
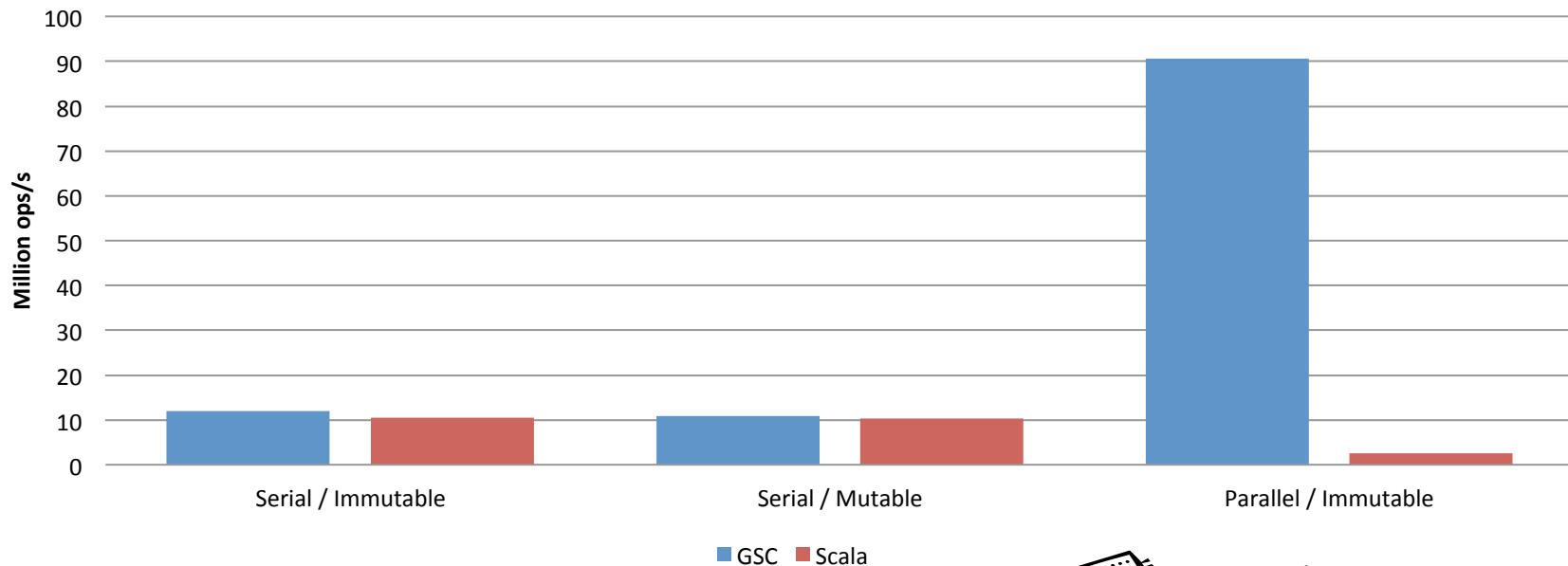
```
this.gscImmutable.asParallel(this.executorService, BATCH_SIZE)
```

Persistent Data Structures

- Assumption: Scala's tree should parallelize moderately well
- Assumption: GSC's array should parallelize very well

Persistent Data Structures

Sorted Set Iteration throughput (higher is better)



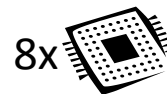
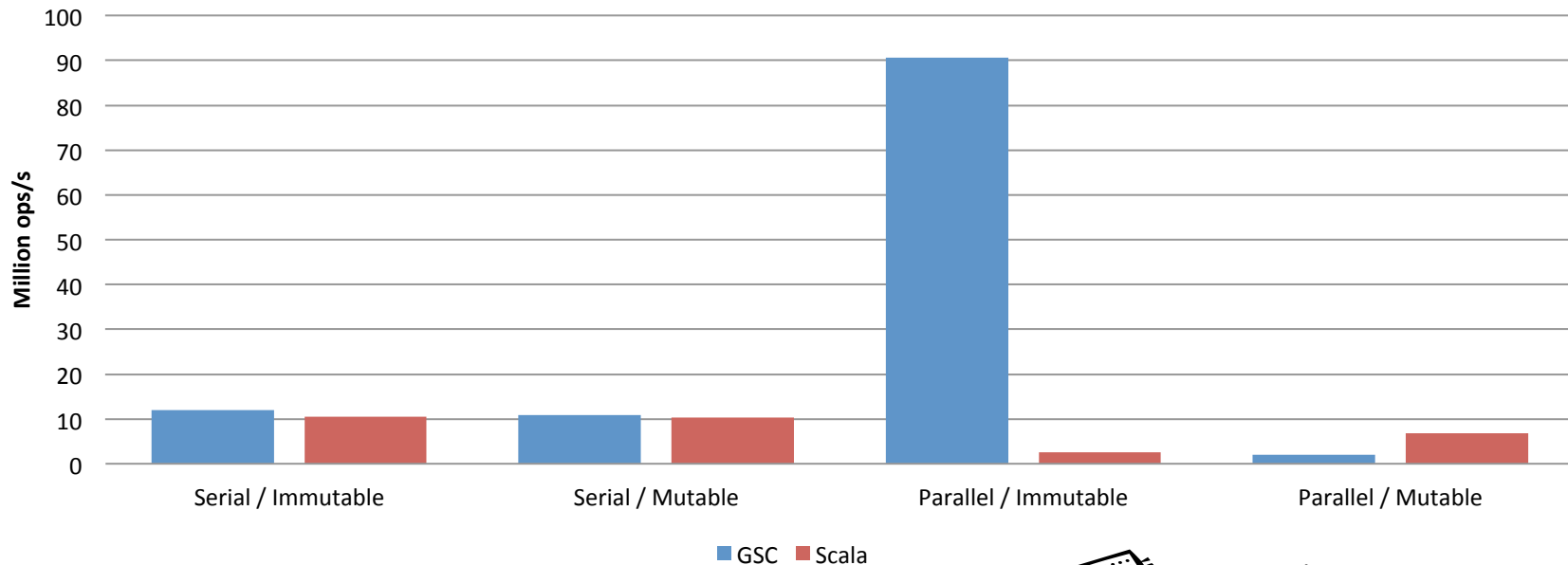
Measured on an 8 core Linux VM
Intel Xeon E5-2697 v2

Persistent Data Structures

- Scala's `immutable.TreeSet` doesn't override `.par`, so parallel is slower than serial
- Some tree operations (like `filter`) are hard to parallelize
- `TreeSet.count()` should be easy to parallelize using fork/join with some work
- New assumption: Mutable trees won't parallelize well either

Persistent Data Structures

Sorted Set Iteration throughput (higher is better)



Measured on an 8 core Linux VM
Intel Xeon E5-2697 v2

Persistent Data Structures

- GS Collections follow up:
 - `TreeSortedSet` delegates to `java.util.TreeSet`. Write our own tree and override `.asParallel()`
- Scala follow up:
 - Override `.par`

Persistent Data Structures

- Proposal: Mutable, Persistent, and **Immutable**
- All three are useful in different cases
- How common is it to share structure, really?

Hash Tables

Hash Tables

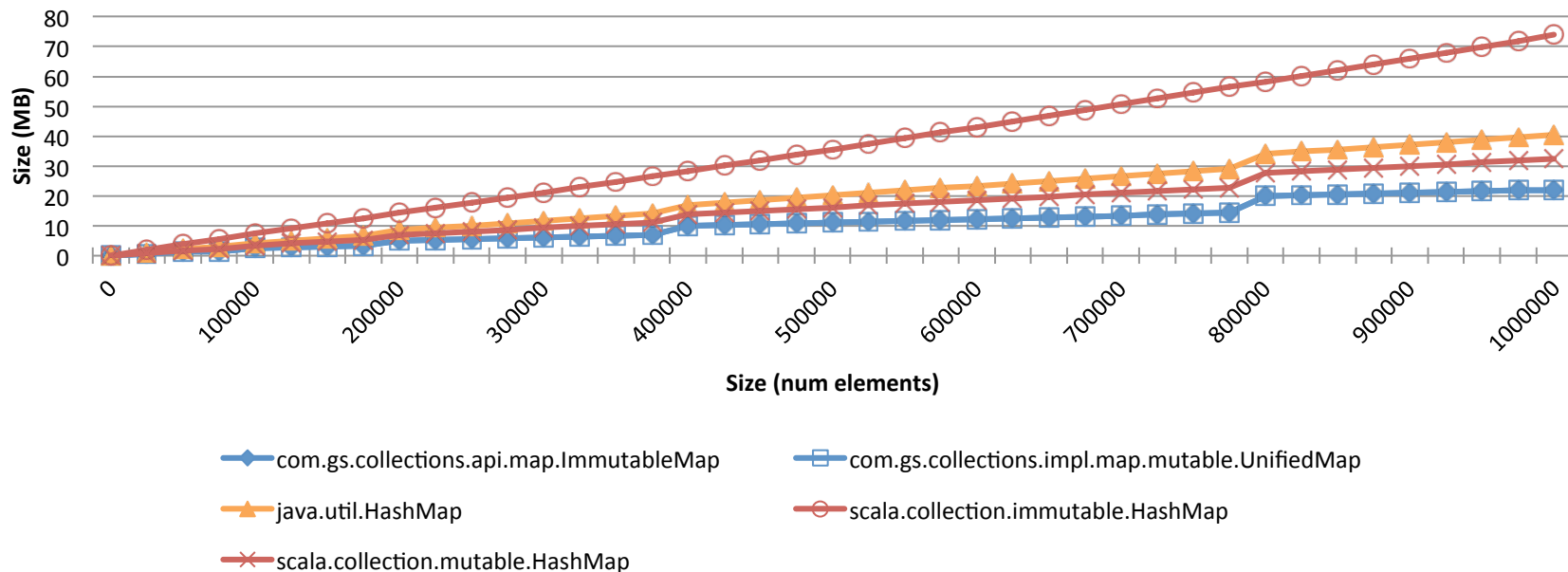
- Scala's `immutable.HashMap` is a hash array mapped trie (persistent structure)
- Wikipedia: “achieves almost hash table-like speed while using memory much more economically”
- Let's see

Hash Tables

- Scala's `mutable.HashMap` is backed by an array of Entry pairs
- `java.util.HashMap.Entry` caches the hashcode
- `UnifiedMap<K, V>` is backed by `Object[]`, flattened, alternating K and V
- `ImmutableUnifiedMap` is backed by `UnifiedMap`

Hash Tables

Maps: Size in MB by number of elements

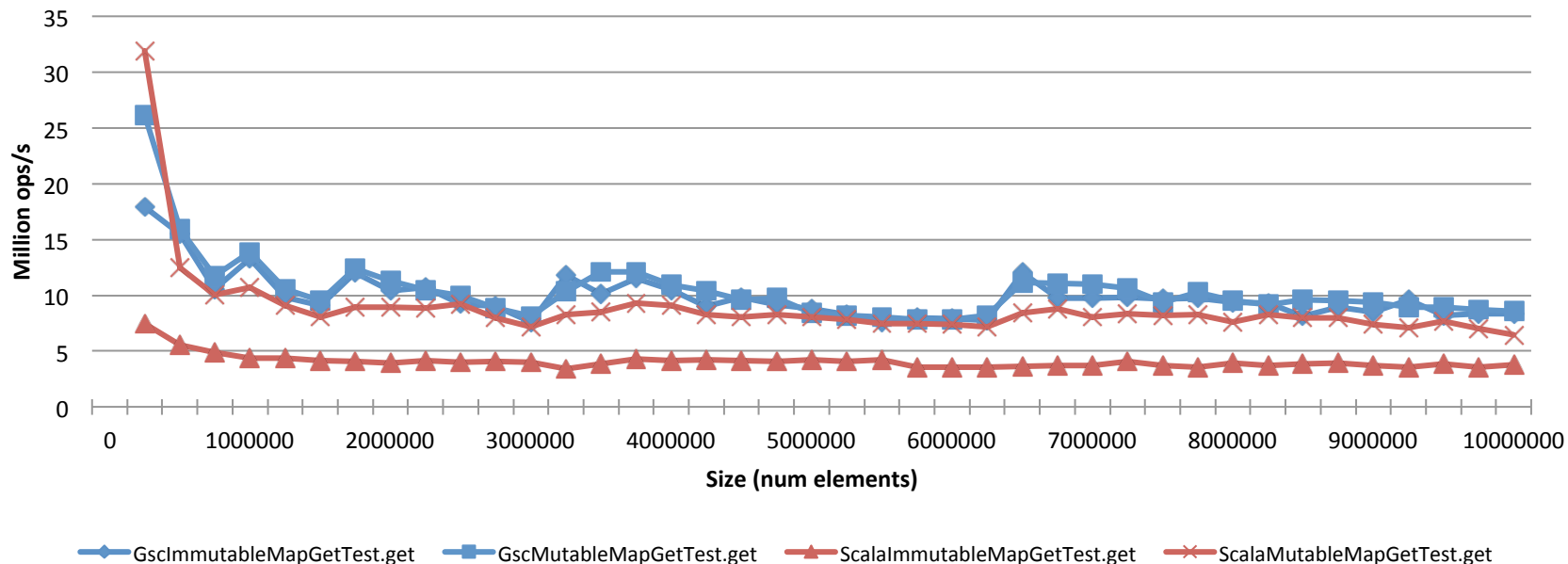


Hash Tables

```
@Benchmark
public void get()
{
    int localSize = this.size;
    String[] localElements = this.elements;
    Map<String, String> localScalaMap = this.scalaMap;
    for (int i = 0; i < localSize; i++)
    {
        if (!localScalaMap.get(localElements[i]).isDefined())
        {
            throw new AssertionError(i);
        }
    }
}
```

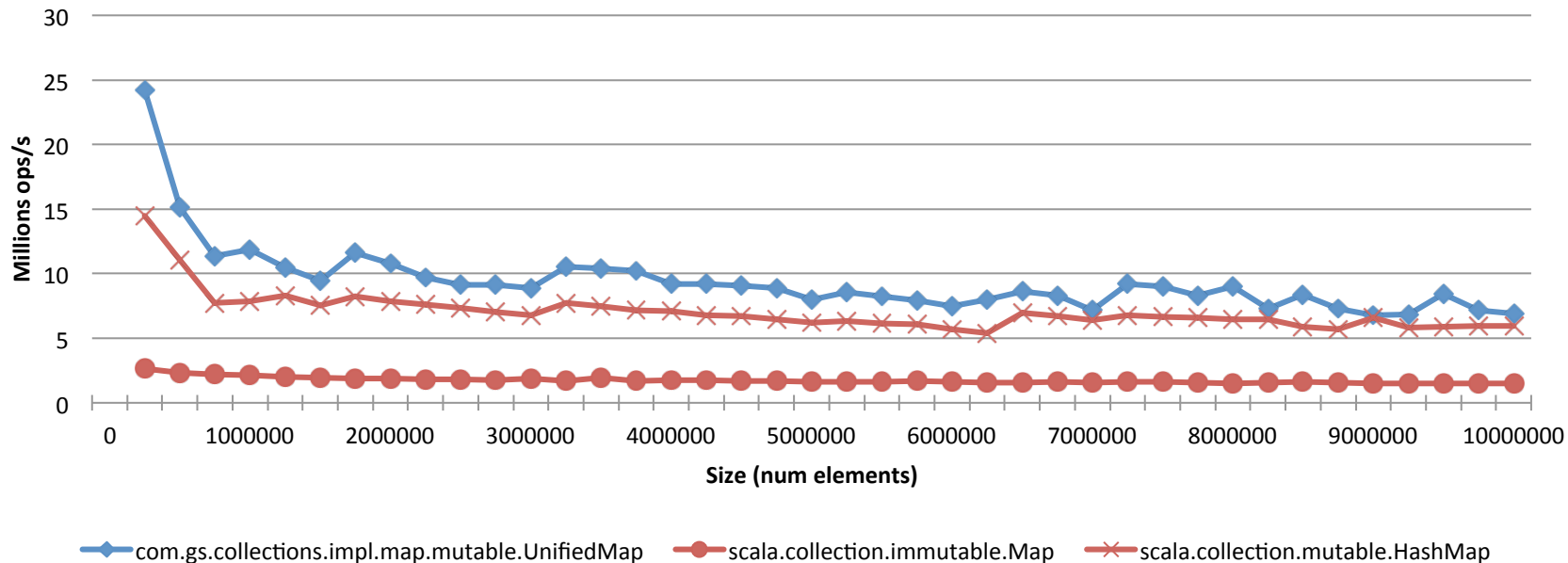
Hash Tables

Map get() throughput by map size (higher is better)



Hash Tables

Map put() throughput by map size (higher is better)



Hash Tables

@Benchmark

```
public mutable.HashMap<String, String> mutableScalaPut()  
{  
    int localSize = this.size;  
    String[] localElements = this.elements;  
  
    mutable.HashMap<String, String> map =  
        new PresizableHashMap<>(localSize);  
  
    for (int i = 0; i < localSize; i++)  
    {  
        map.put(localElements[i], "dummy");  
    }  
    return map;  
}
```

Hash Tables

HashMap ought to have a constructor for pre-sizing

```
class PresizableHashMap[K, V](val _initialSize: Int)
  extends scala.collection.mutable.HashMap[K, V]
{
  private def initialCapacity =
    if (_initialSize == 0) 1
    else smallestPowerOfTwoGreaterThan(
      (_initialSize.toLong * 1000 / _loadFactor).asInstanceOf[Int])

  private def smallestPowerOfTwoGreaterThan(n: Int): Int =
    if (n > 1) Integer.highestOneBit(n - 1) << 1 else 1

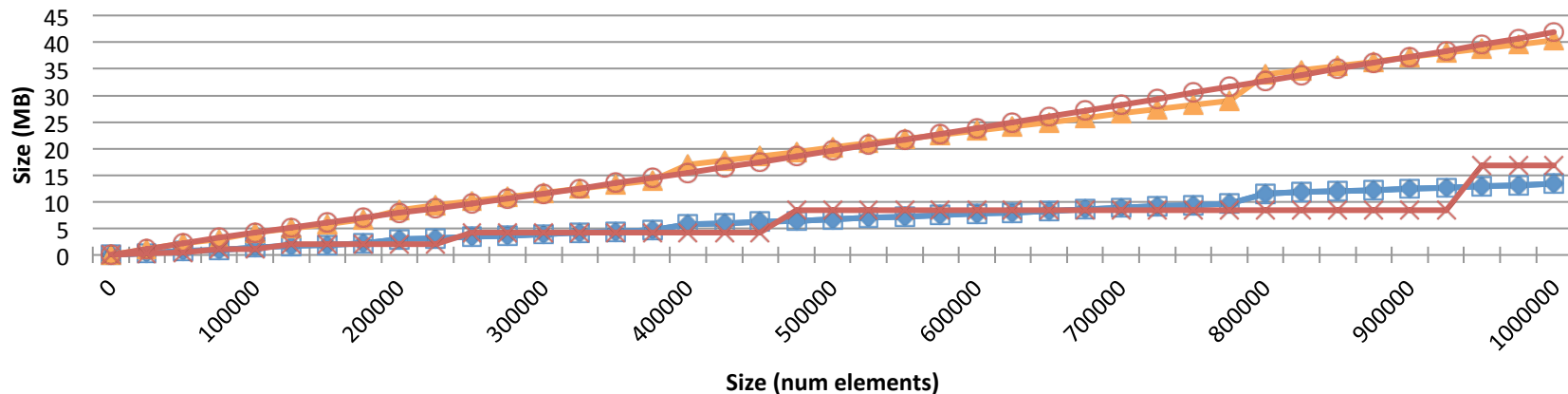
  table = new Array(initialCapacity)
  threshold = ((initialCapacity.toLong * _loadFactor) / 1000).toInt
}
```


Hash Tables

- `scala.collection.mutable.HashSet` is backed by an array using open-addressing
- `java.util.HashSet` is implemented by delegating to a `HashMap`.
- `UnifiedSet` is backed by `Object[]`, either elements or arrays of collisions
- `ImmutableUnifiedSet` is backed by `UnifiedSet`

Hash Tables

Sets: Size in MB by number of elements



com.gs.collections.api.set.ImmutableSet

com.gs.collections.impl.set.mutable.UnifiedSet

java.util.HashSet

scala.collection.immutable.HashSet

scala.collection.mutable.HashSet

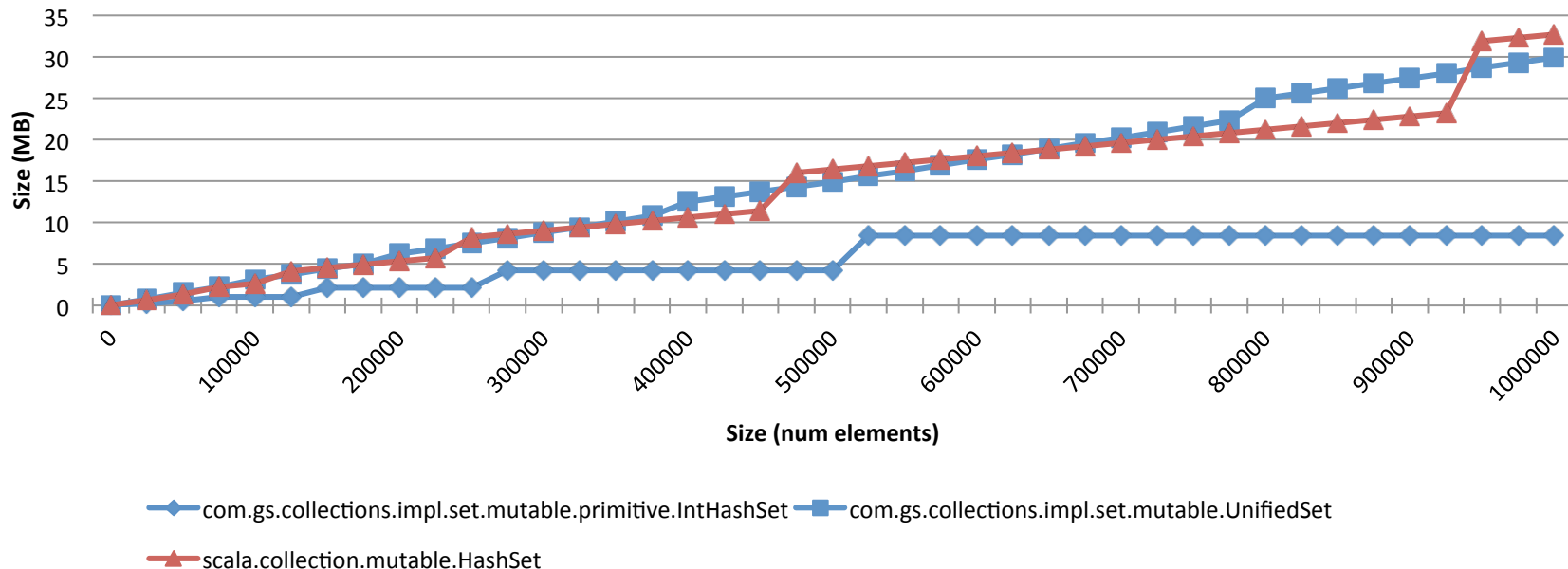
Primitive Specialization

Primitive Specialization

- Boxing is expensive
 - Reference + Header + alignment
- Scala has specialization, but most of the collections are not specialized
- If you cannot afford wrappers you can:
 - Use primitive arrays (only for lists)
 - Use a Java Collections library

Primitive Specialization

Sets: Size in MB by number of elements



Primitive Specialization

- Proposal: Primitive lists, sets, and maps in Scala
- Not Traversable – outside the collections hierarchy
- Fix Specialization on Functions (lambdas)
 - Want for-comprehensions to work well

Fork/Join


Parallel / Lazy / Scala

```
val list = this.integers
  .par
  .filter(each => each % 10000 != 0)
  .map(String.valueOf)
  .map(Integer.valueOf)
  .filter(each => (each + 1) % 10000 != 0)
  .toBuffer
```

```
Assert.assertEquals(999800, list.size)
```



Parallel / Lazy / GSC

```
MutableList<Integer> list = this.integersGSC   
    .asParallel(this.executorService, BATCH_SIZE)  
    .select(each -> each % 10_000 != 0)  
    .collect(String::valueOf)  
    .collect(Integer::valueOf)  
    .select(each -> (each + 1) % 10_000 != 0)  
    .toList();
```

```
Verify.assertSize(999_800, list);
```

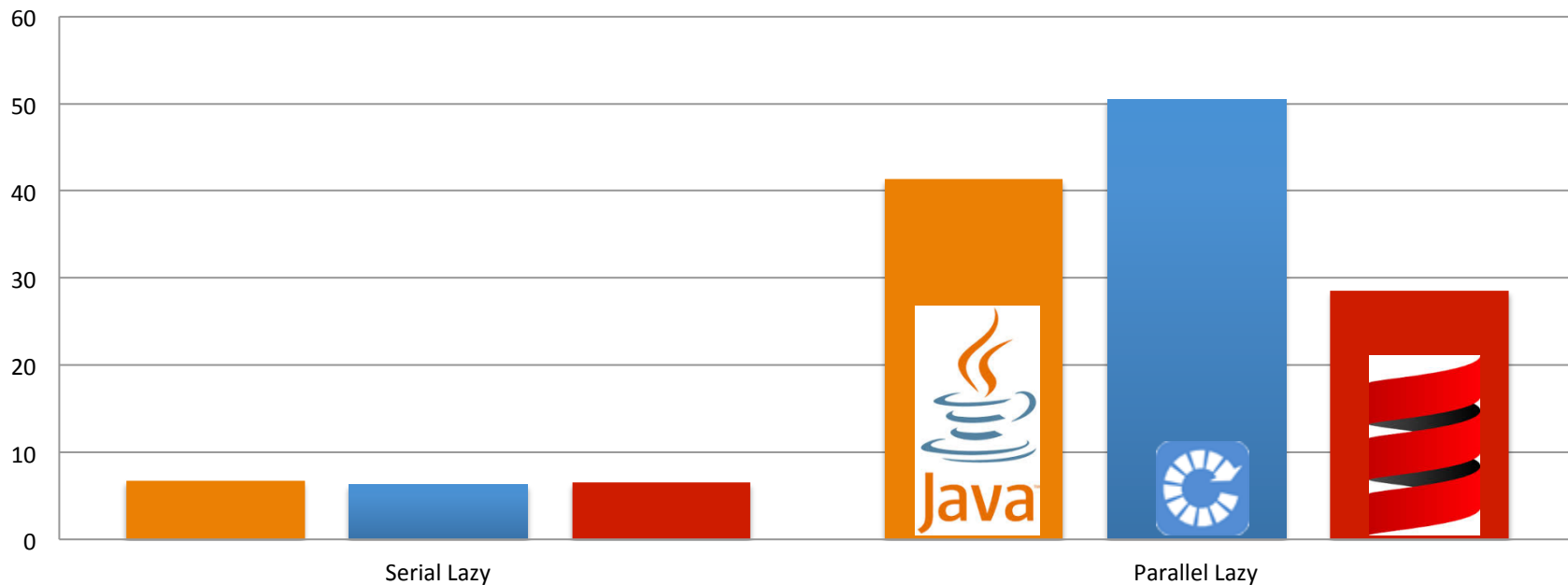
Parallel / Lazy / JDK

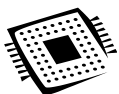
```
List<Integer> list = this.integersJDK
    .parallelStream()
    .filter(each -> each % 10_000 != 0)
    .map(String::valueOf)
    .map(Integer::valueOf)
    .filter(each -> (each + 1) % 10_000 != 0)
    .collect(Collectors.toList());
```

```
Verify.assertSize(999_800, list);
```



Stacked computation ops/s (higher is better)

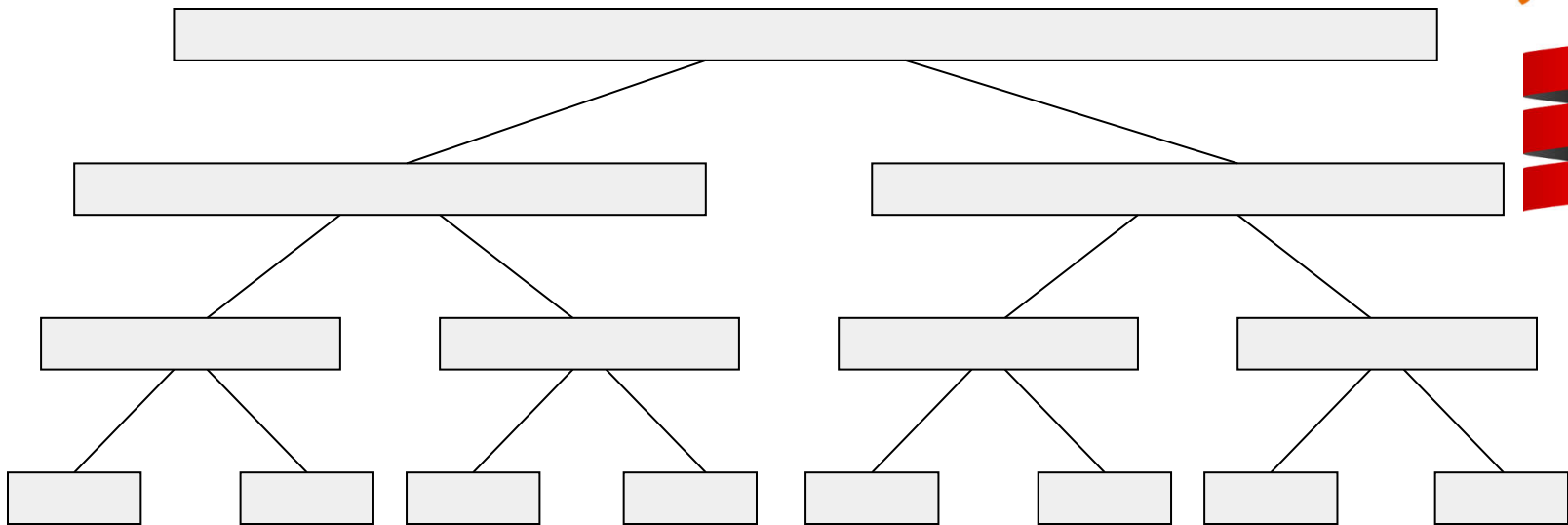


8x 

■ Java 8 ■ GS Collections ■ Scala

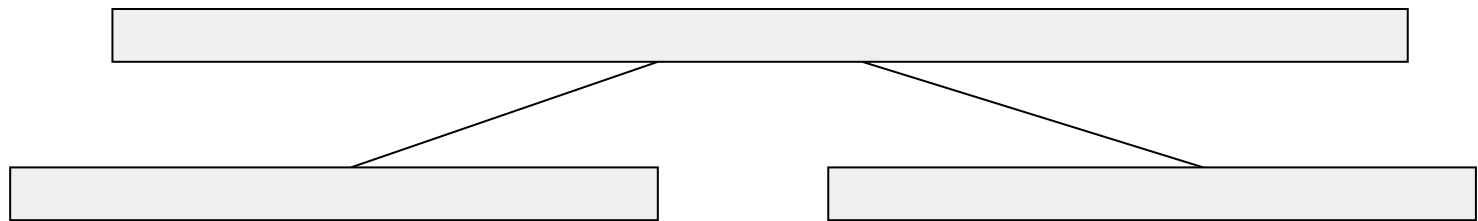
Fork-Join Merge

- Intermediate results are merged in a tree
- Merging is $O(n \log n)$ work and garbage



Fork-Join Merge

- Amount of work done by last thread is $O(n)$



Parallel / Lazy / GSC

```
MutableList<Integer> list = this.integersGSC  
    .asParallel(this.executorService)  
    .select(each -> each % 10_000)  
    .collect(String::valueOf)  
    .collect(Integer::valueOf)  
    .select(each -> (each + 1) % 10_000 != 0)  
    .toList();
```



ParallelIterable.toList() returns a CompositeFastList, a List with O(1) implementation of addAll()

```
Verify.assertSize(999_800, list);
```

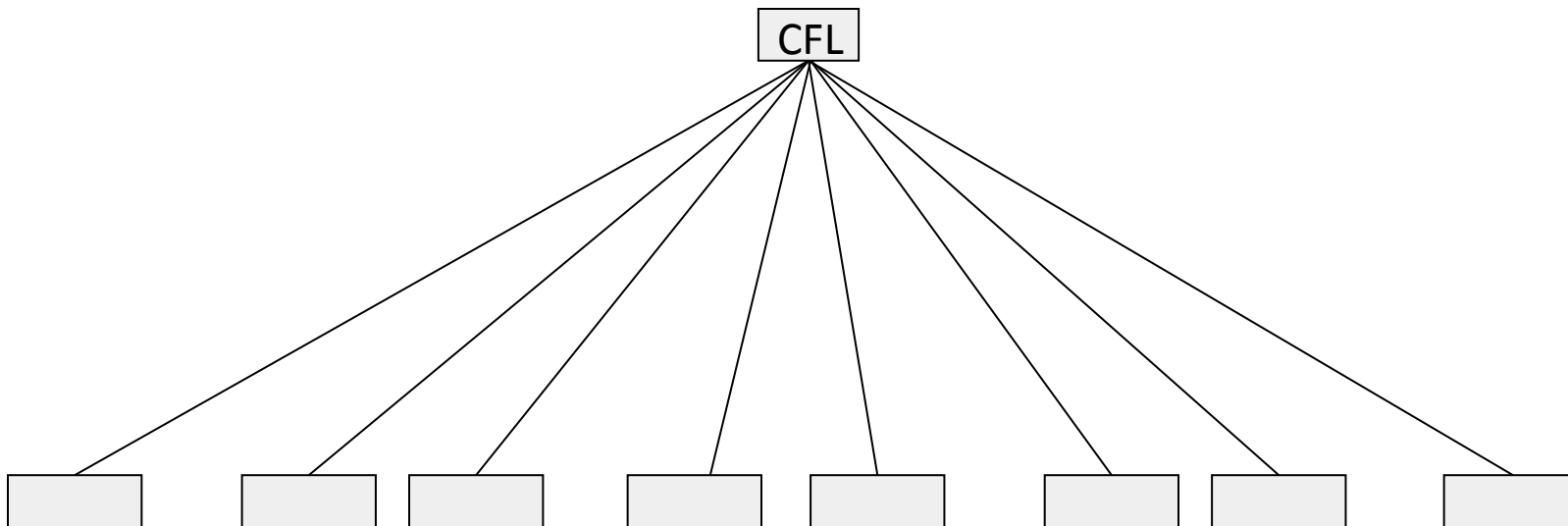
Parallel / Lazy / GSC



```
public final class CompositeFastList<E> {  
    private final FastList<FastList<E>> lists =  
        FastList.newList();  
  
    public boolean addAll(Collection<? extends E> collection) {  
        FastList<E> collectionToAdd = collection instanceof FastList  
            ? (FastList<E>) collection  
            : new FastList<E>(collection);  
        this.lists.add(collectionToAdd);  
        return true;  
    }  
    ...  
}
```

CompositeFastList Merge

- Merging is $O(1)$ work per batch



Fork/Join

- Fork-join is general purpose but always requires merge work
- We can get better performance through specialized data structures meant for combining

Summary

GS Collections and Scala's Collections are different

- Persistent data structures
- Hash tables
- Primitive specialization
- Fork-join

More Information

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@GoldmanSachs

stackoverflow.com/questions/tagged/gs-collections

github.com/goldmansachs/gs-collections/tree/master/jmh-tests

github.com/goldmansachs/gs-collections/tree/master/memory-tests

github.com/goldmansachs/gs-collections-kata

infoq.com/presentations/java-streams-scala-parallel-collections



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BUILD

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Q&A